

Ring Signatures and Other Cryptographic Research Challenges in Blockchain

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About Me



Yuen Tsz Hon (John)
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- Assistant Professor (2018 – present)
 - Department of Computer Science + Faculty of Economics
 - Research area: Fintech, cryptography



- Senior Researcher (2013 – 2018)
Research area: blockchain, big data privacy, searchable encryption for database, identity-based authentication for 5G, etc. Published 14 papers including IEEE TC, ESORICS, and filed 6 patents.



- Post-doctoral Fellow (2010 – 2013)
 - Research on cryptography and security. Published 16 papers including Eurocrypt, ESORICS.



Senior Research
Associate (2010)

UNIVERSITY OF
WOLLONGONG
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Ph.D. (2007-2010)



M.Phil. (2004-2006)
B.Eng. (2001-2004)

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02 Blockchain Research Challenges

03 Ring Signatures in Blockchain

Blockchain: The Trust Machine

➤ Blockchain: underlying technology of Bitcoin

“The blockchain lets people who have no particular confidence in each other collaborate without having to go through a neutral central authority. Simply put, it is a **machine for creating trust.**”

Economist, October 2015



Why do we need Blockchain?

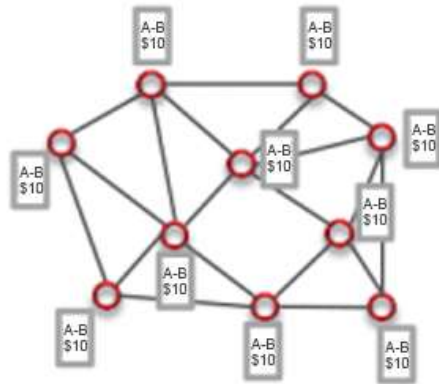
Goal: Transfer value digitally without trusted third party

Problem:

Resolve conflicts



Distributed Ledger



- Single version of history
- Transparent
- Immutable
- Unforgeable (signature)

Blockchain



- Order transactions

Consensus



- Determine who can record



: Hash function

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Challenges of Using Blockchain in FinTech

Business obstacles

Cryptocurrency

Bubble



- Promising applications are overwhelmed by unrealistic projects due to speculative investments

Auditing



- Difficult to find out the real user identity under pseudonym

Regulation



- Lack of law for cryptocurrency, smart contract and Fintech use cases (e.g. ICO, P2P lending, etc.)

Technical obstacles

Efficiency



- Bitcoin: 7 transactions per second (tps)
- Most blockchain system: <10,000 tps
 - Cf. Visa with peak capacity 56,000 tps, Alibaba with max. 256,000 tps on 2017.11.11

Privacy



- Public transaction history → no privacy of transaction amount/address
- Use cryptography → Tradeoff between privacy and efficiency

Security



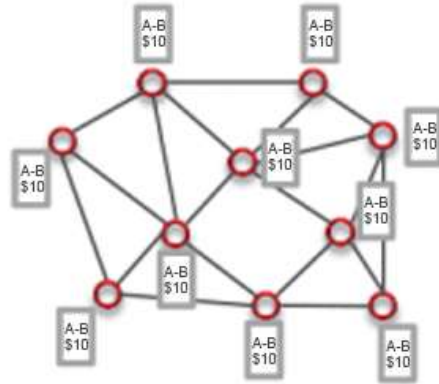
- Software security of user's wallet and exchange
 - Major exchange Mt. Gox was bankrupted after USD\$450M bitcoin was stolen

Technical Challenge 1: Consensus

Problem:
Resolve conflicts



Distributed Ledger



Blockchain



Consensus

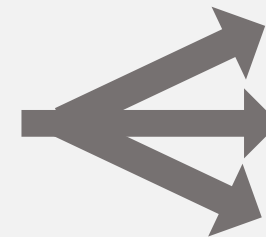


➤ Determine who
can record



Problem of Bitcoin's Consensus (Proof of Work):

- Low throughput (7 tps)
- High latency (10 minutes/block, 6 blocks for finality)
- Waste electricity



Theoretical breakthrough

Semi-trusted consensus nodes

Trusted hardware

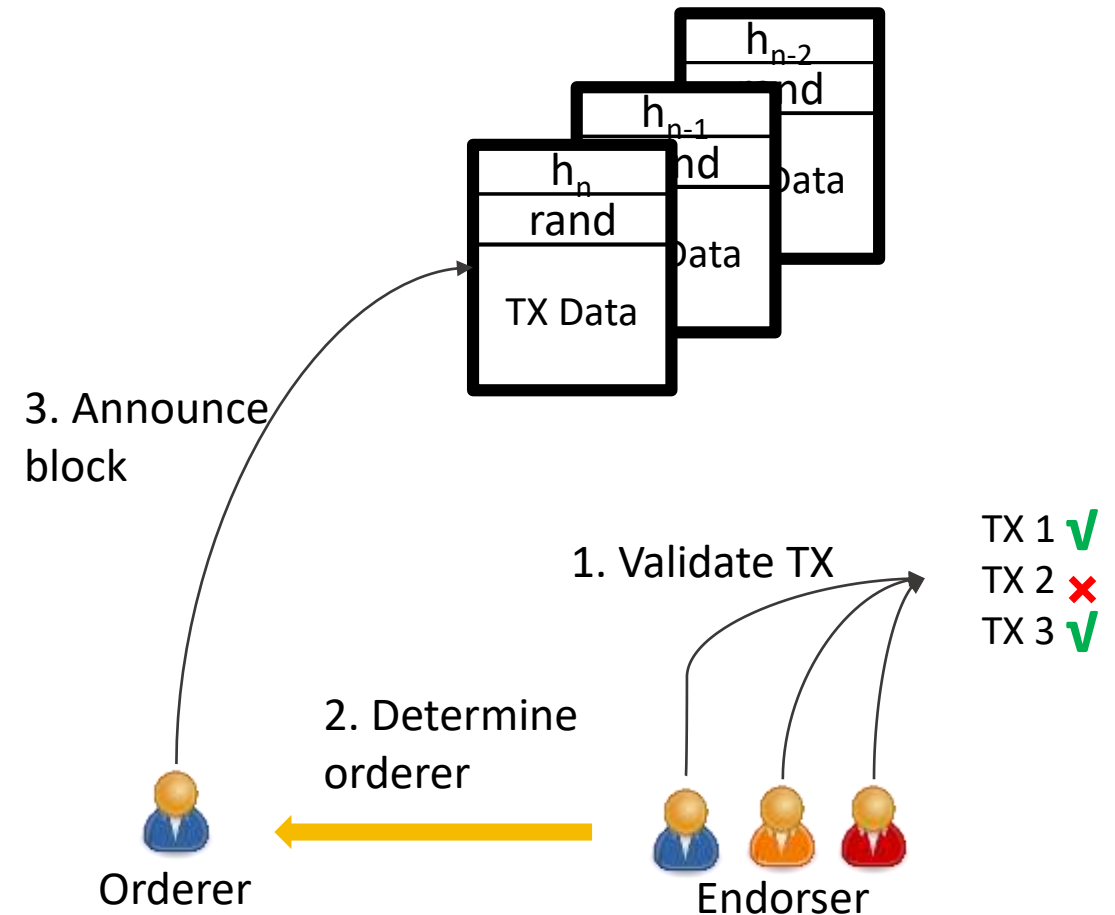
Technical Challenge 1: Consensus (1)

Theoretical breakthrough: New consensus algorithms using cryptographic primitives, such as threshold signature, secret sharing, etc.

- Proof of Stake (PoS)
 - Creator of the next block is chosen via combinations of random selection and wealth (stake)
 - Used in various cryptocurrencies such as Ethereum

Active research area:

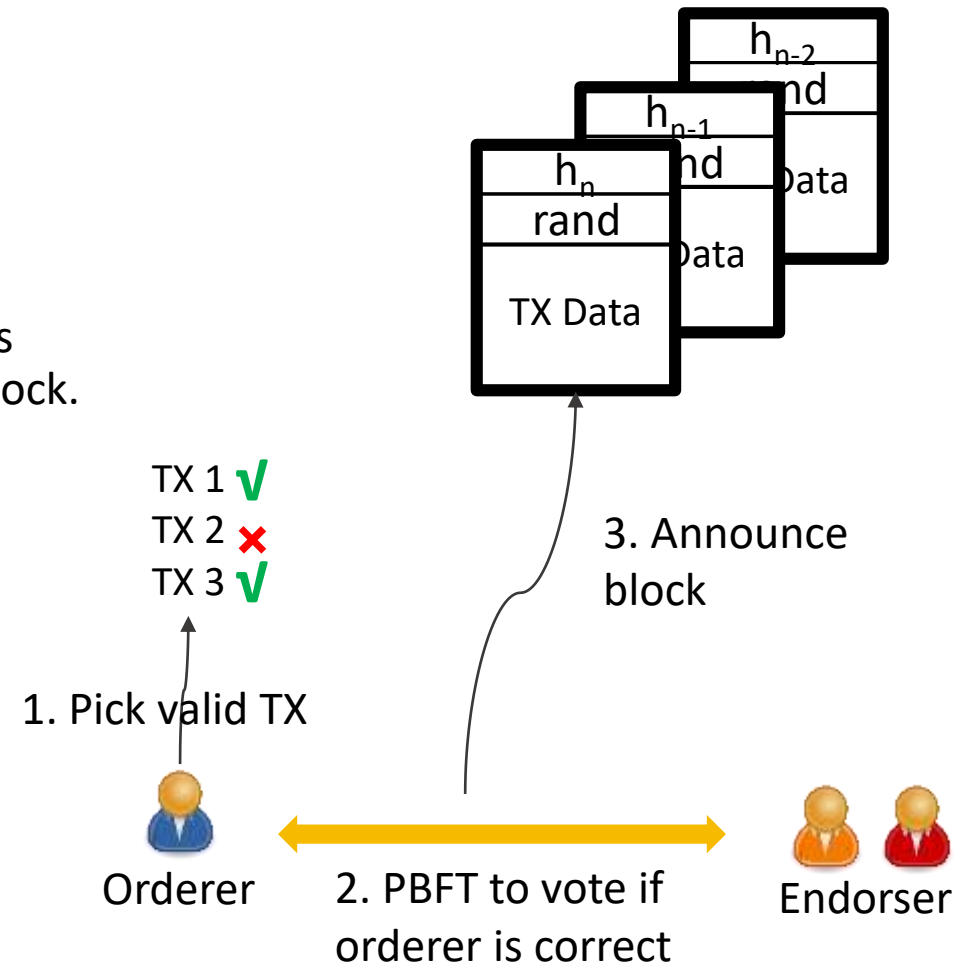
- Honey Badger (Miller et al. CCS 2016)
- Ouroboros (Kiayias et al. CRYPTO 2017)
 - used in Cardano (#5 cryptocurrency, USD\$16 billion)
 - Ouroboros praos (Eurocrypt 2018), Ouroboros genesis
- Algorand (Micali et al. SOSP 2017), Algorand agreement
- Snow White (Pass, Shi. Asiacrypt 2017), Thunderella (Eurocrypt 2018)



Technical Challenge 1: Consensus (2)

Semi-trusted consensus nodes: Assume consensus nodes are known and majority are honest, it can be reduced to Byzantine fault tolerance algorithm.

- Practical Byzantine Fault Tolerance (PBFT)
 - What is PBFT:
 - Achieve consensus if $>2/3$ nodes are not faulty
 - How to use PBFT in blockchain:
 - Nodes become orderer in round-robin. All nodes “vote” if they agree with the orderer’s written block.
 - Advantages:
 - Efficient: Theoretically 10000+ tps
 - Well-studied algorithm
 - Disadvantages:
 - Unscalable: Only allows < 20 nodes
 - Nodes must be known to each other



Technical Challenge 1: Consensus (2)

Semi-trusted consensus nodes: More suitable to Fintech applications

Consortium Blockchain:

- Blockchain with multi-authorities
- More realistic for Fintech applications
 - E.g. Clearing and settlement between banks
- Industrial alliances (e.g. R3, Hyperledger, Enterprise Ethereum Alliance) are formed
- PBFT-type consensus is suitable
 - Good for low-medium frequency trading (10,000 tps)
 - Daily transaction??
 - Visa with peak capacity 56,000 tps
 - Alibaba with 325,000 tps at peak on 2017.11.11

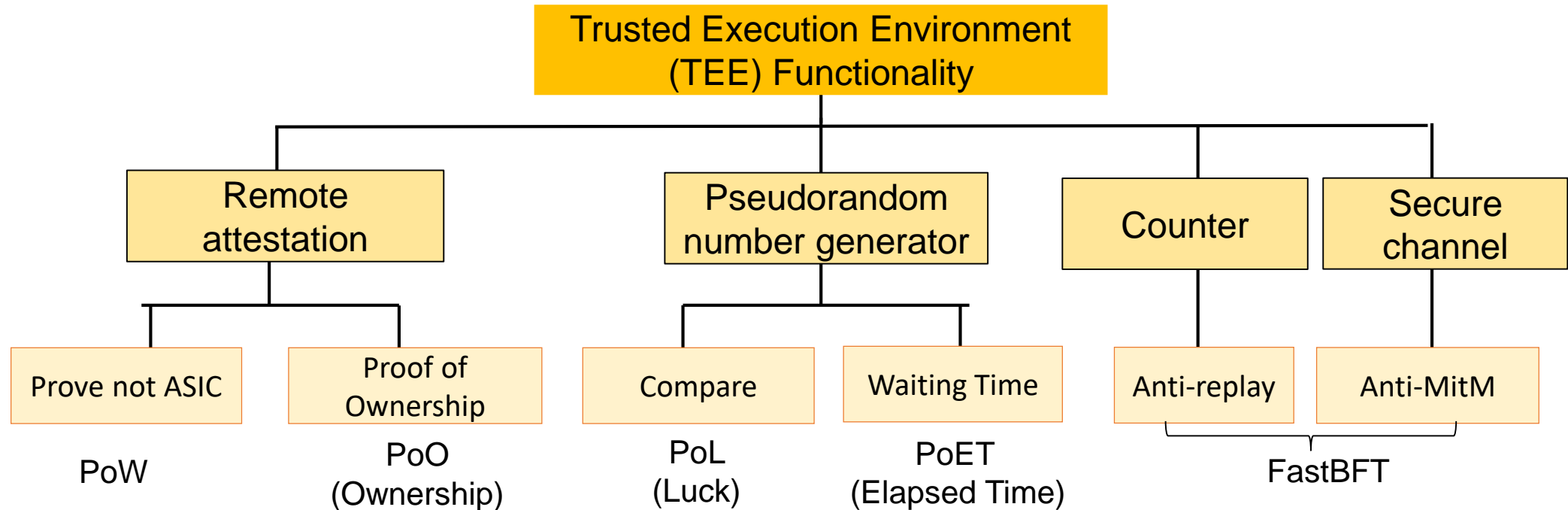


- Open source blockchain alliance to advance cross-industry collaboration



Technical Challenge 1: Consensus (3)

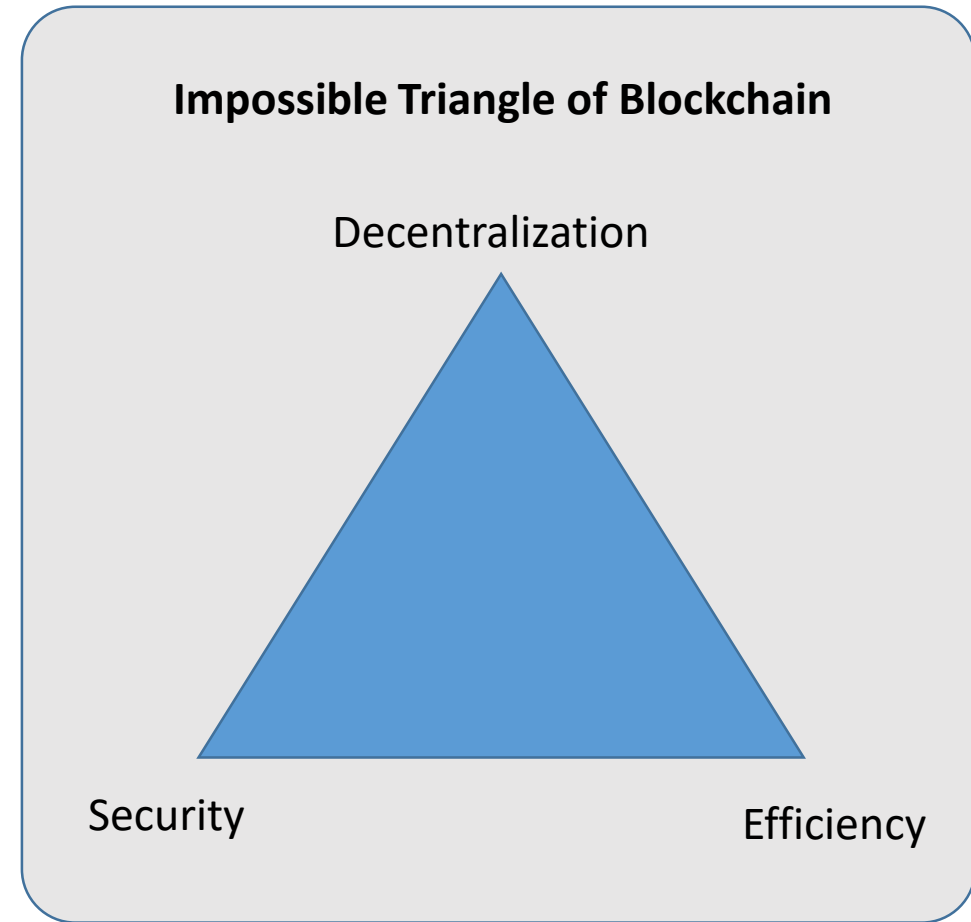
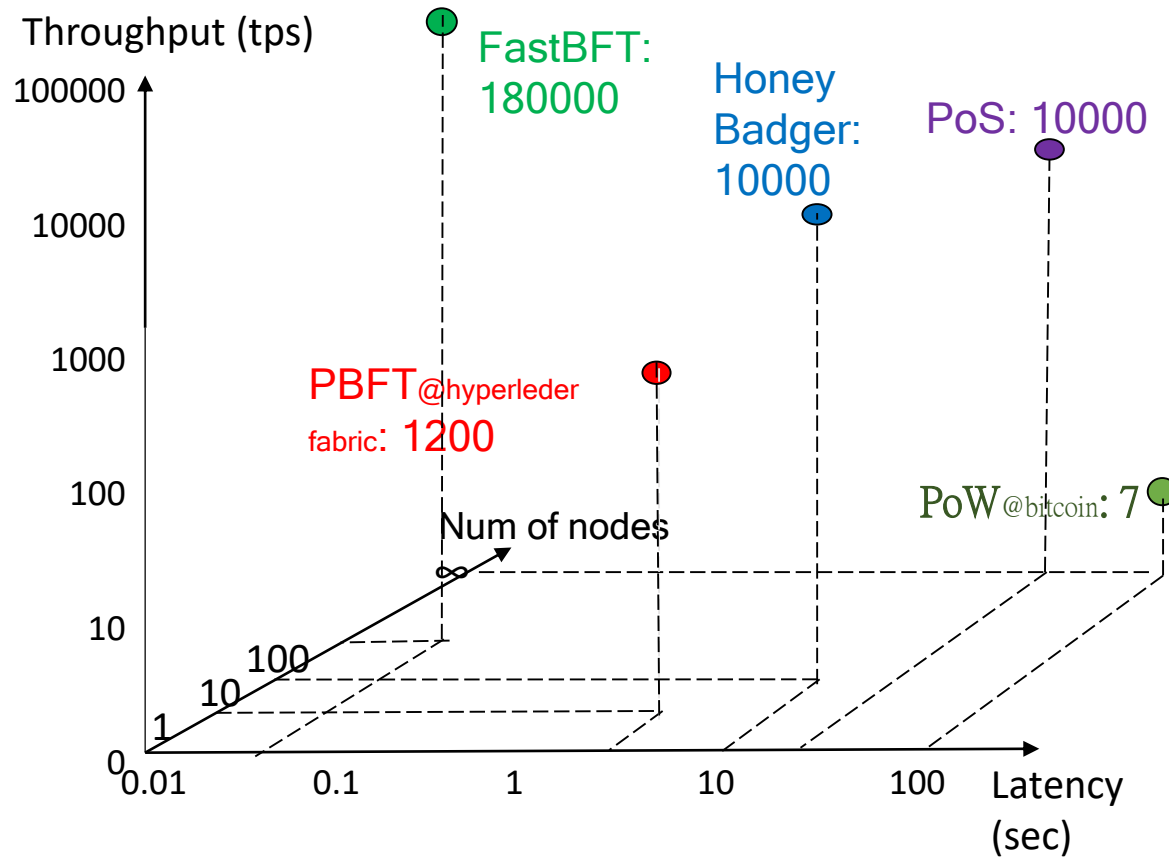
Trusted hardware: Most efficient solution based on strong security requirement



Hardware-based consensus algorithm is more concerned by the industry (e.g. Intel's PoET, NEC's FastBFT).

Technical Challenge 1: Consensus Summary

➤ Tradeoff:



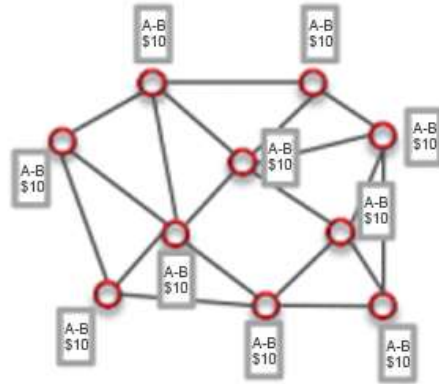
Technical Challenge 2: Data Structure

Problem:

Resolve conflicts



Distributed Ledger



Blockchain



➤ Order transactions

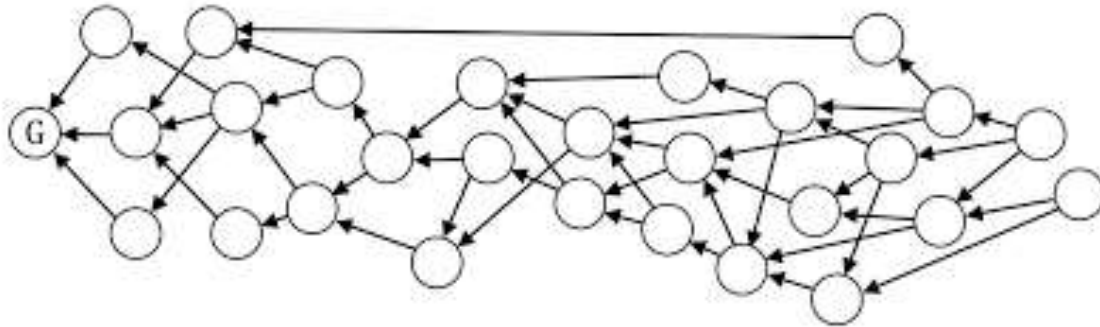
Consensus



- Why do we need a chain of blocks? ➡ In case of conflict, discard the second transaction
- Problem of “blockchain” data structure: no parallel computation → limited throughput

Technical Challenge 2: Data Structure

Directed acyclic graph (DAG)



- DAG: Each transaction references two or more previous transactions → parallel transactions are possible
- Question: How to solve conflicts?



IOTA:

- Support currency transaction only
- All nodes are weighted. Heavier branches are more likely to survive
- Markov Chain Monte Carlo is used to choose which branch to extend



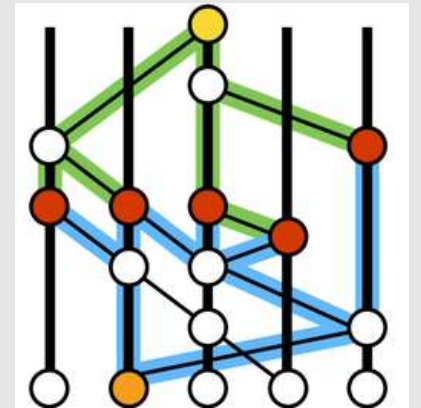
Byteball:

- Support smart contract
- Trusted *witness* is responsible to finalize certain blocks (chain-then-block)



Hashgraph:

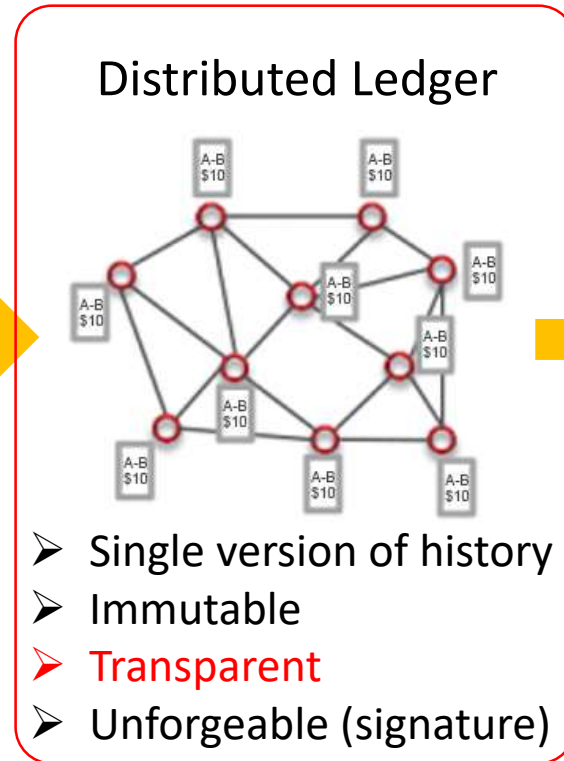
- DAG for permissioned blockchain



- Also Graphchain by boyen et. al. in BCC@AsiaCCS 2018

Technical Challenge 3: Privacy

Problem:
Resolve conflicts



Blockchain



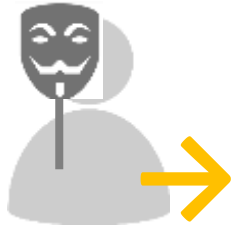
Consensus



- Transparency is good for some use cases, e.g. supply chain management
- Transparency is not desirable for sensitive information, e.g. transaction details between two banks should not be known to other banks

Technical Challenge 3: Privacy

3 Types of Privacy:



Sender
anonymity



Confidential
transaction



Recipient
anonymity



Zcash (#26)

- Use zero knowledge proof of circuit “zk-SNARK” to achieve all 3 types of privacy
- Based on IEEE SP 2014 paper

✓ Can be extended to zk-smart contract (IEEE SP 2016)

- ✗ 1. Prove time = 30 sec.
- 2. Require trusted setup



Monero (#13)

- Sender anonymity: linkable ring signature
- Recipient anonymity: Diffie-Hellman key exchange
- Confidential transaction: Pederson commitment

✓ Quite Practical

- ✗ Not scalable for sender anonymity



Dash (#12)

- Mixing a few transactions' input and output together, using homomorphic encryption

✓ Practical

- ✗ Limited anonymity

Technical Challenge 3: Privacy

Privacy: Popular among both researchers and practitioners



Zero-knowledge proof approach:

- zk-SNARK is implemented on top of Ethereum (#2) and JP Morgan's Quorum platform
- Very active research area, 10+ papers in top conferences since 2014



Coin-shuffling approach:

- Improved in applications, such as SharedCoins, Dark Wallet.
- CoinShuffle paper in ESORICS 2014, works for bitcoin



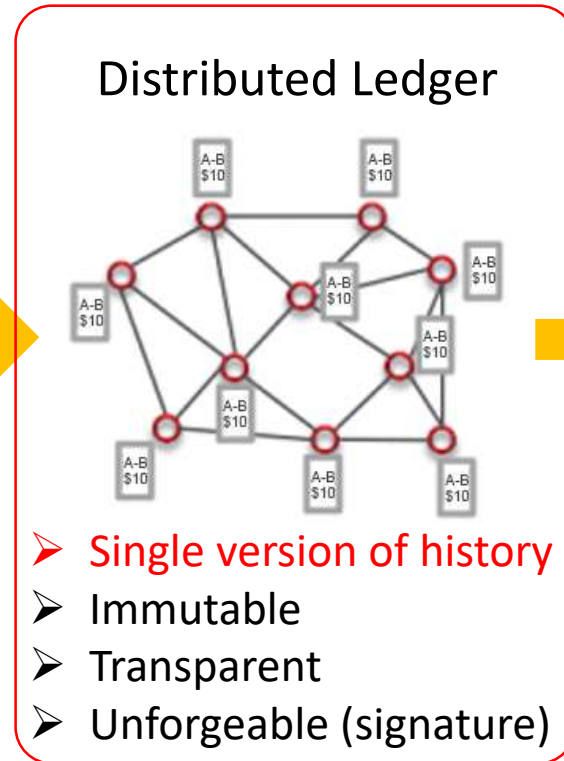
Linkable Ring Signature + Confidential Transaction approach:

- Confidential transaction is used by top startups such as Blockstream and Chain.com
- RuffCT is proposed by Monero developers to improve the scalability of linkable ring signature

Technical Challenge 4: Sharding

Problem:

Resolve conflicts



Blockchain



Consensus



- Sharding is a common technique in database. It improves the throughput of blockchain system
 - Network and Transaction Sharding: Each subgroup of nodes reach consensus on a subset of transactions in parallel, but every node will need to store all the data. (e.g., Zilliqa, CCS 2016)
 - State sharding: State data is split and stored on different shards separately. Before the sharding group is re-organized, data exchanges must be done in advance. (e.g. Ethereum 2.0)

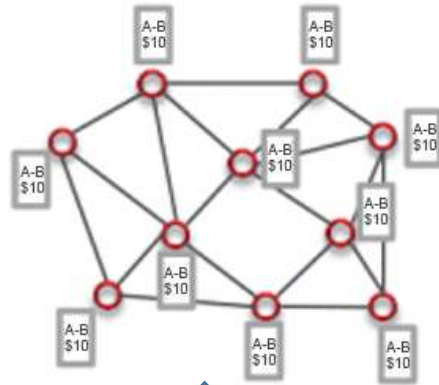
Technical Challenge 5: Layer 2-solution

Problem:

Resolve conflicts



Distributed Ledger



Blockchain



Consensus



Layer 2-solution: Handling transaction off main chain

- **Lightning network:** increasing transaction throughput and confirmation speed of bitcoin network. Submit a fraud proof against the malicious side to the main chain so as to confiscate his deposit as a penalty.
- **State channel:** a more general 2 party off-chain channel which manages the states of smart contract
- **Plasma:** multiple parties can participate and the confirmation time of a transaction depends on when the plasma block header is submitted to the plasma contract on the parent blockchain.
- Truebit, Plasma cash, etc...

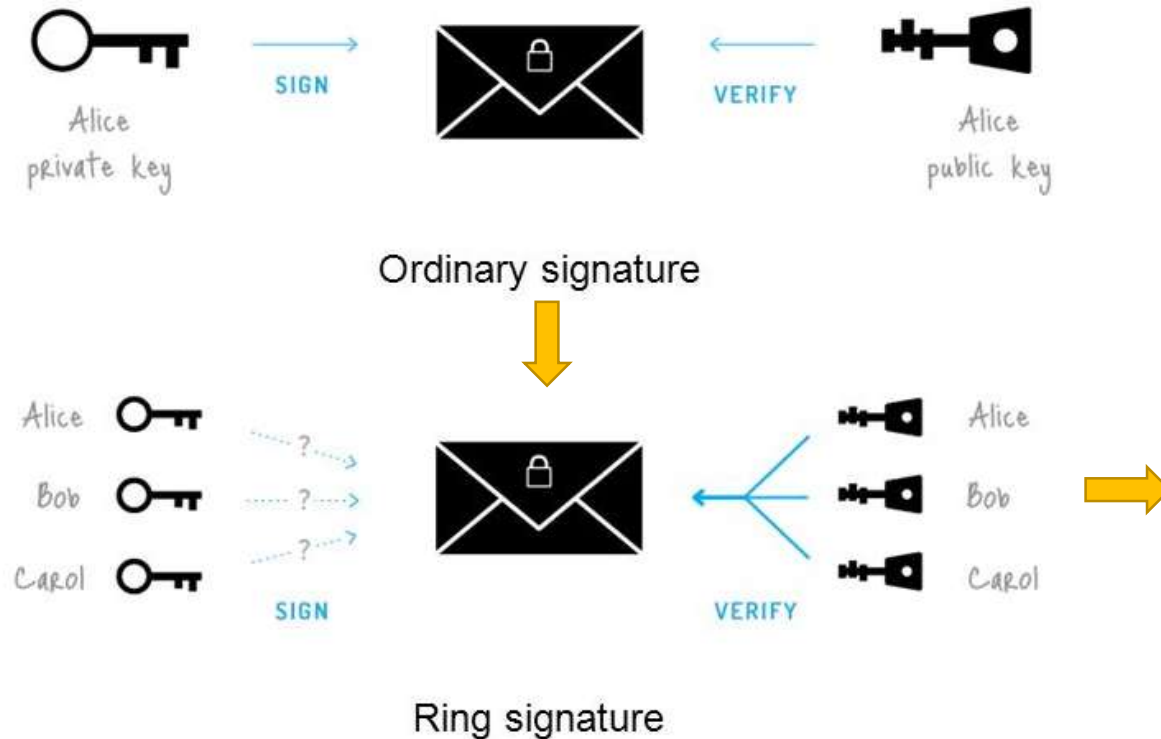
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Ring Signatures for Blockchain Privacy



Linkable Ring Signature:

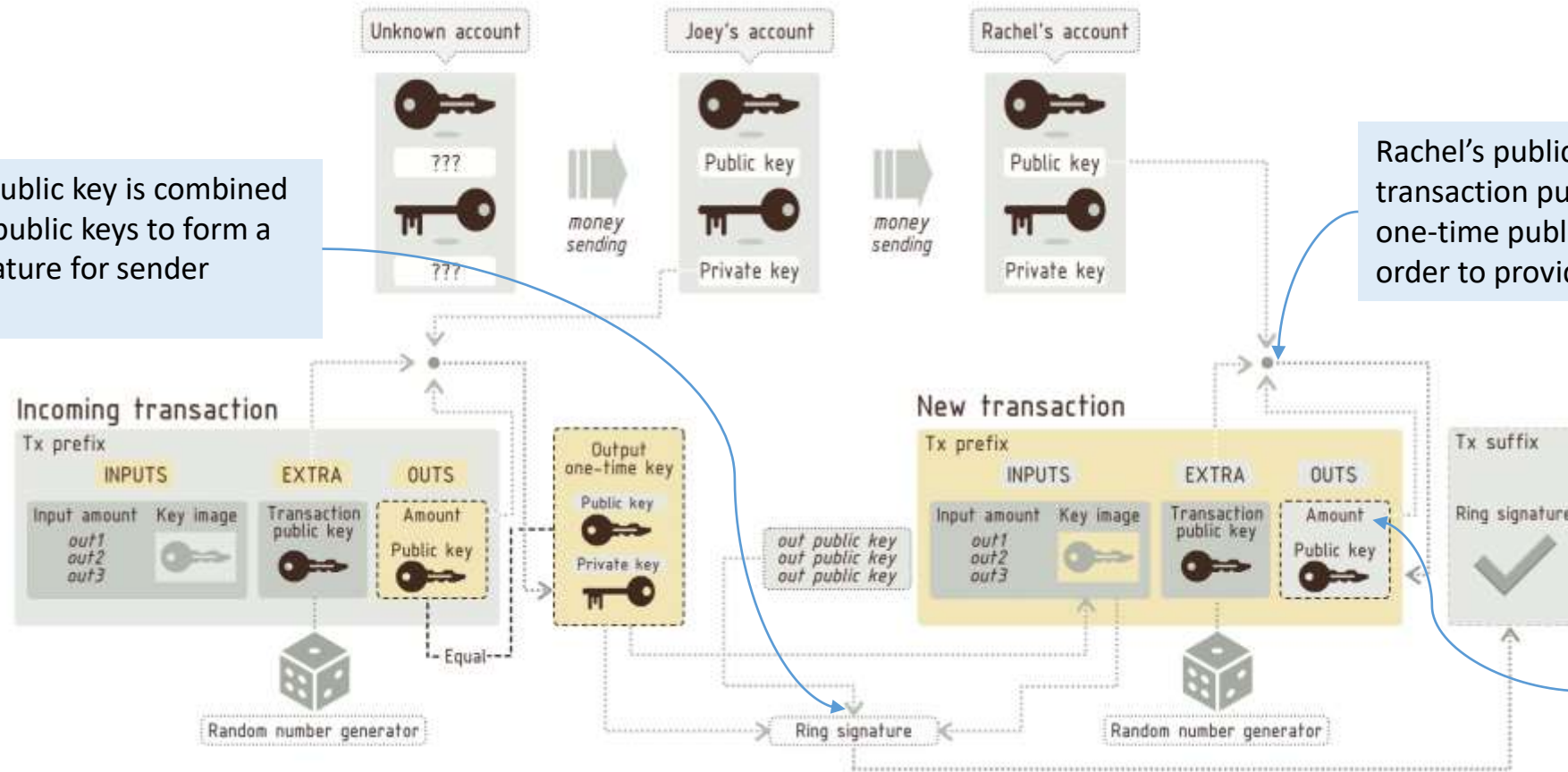
- Detect signer that sign twice → used to detect double spending in blockchain

Ring Signatures for Blockchain Privacy



Confidential Transactions in Monero: Joey receives money from some unknown account (left hand side) and sends it to Rachel (right hand side).

Joey's one-time public key is combined with 3 other out public keys to form a linkable ring signature for sender anonymity.



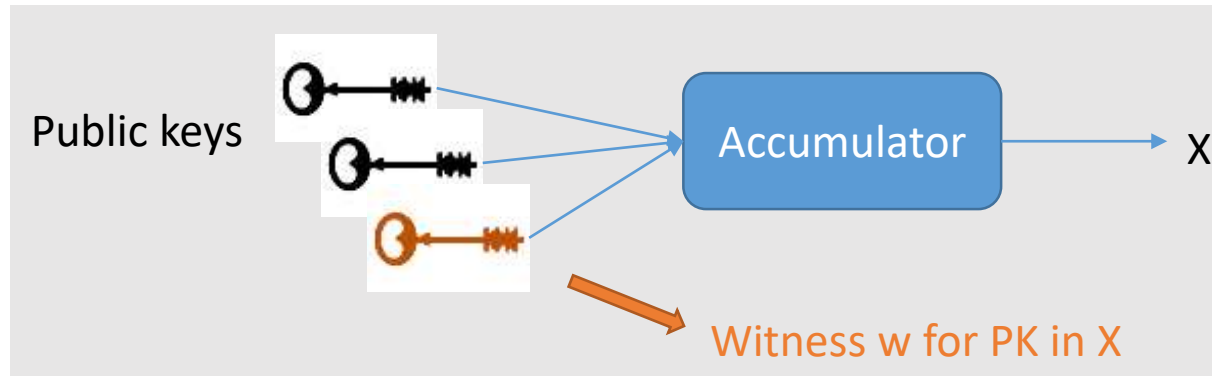
Transaction privacy can be achieved by computing additive homomorphic commitment on the input and output amount.

Monero's limitation:

- Signature size is $O(n)$, where n = number of public keys (PKs) to "hide" the signer → not scalable

RingCT 2.0

- RingCT 2.0: A Compact Accumulator-Based (Linkable Ring Signature) Protocol for Blockchain Cryptocurrency Monero. Shi-Feng Sun, Man Ho Au, Joseph K. Liu, Tsz Hon Yuen. ESORICS 2017.



Linkable Ring Signatures:

PoK{ (SK, PK, w) : SK is the secret key of PK
and w is a witness for PK in X
and Key-Image = F(SK) }

Challenge:

- RingCT protocol is essentially a multi-ring signature scheme
 - If the transaction has 3 inputs, it has a ring signature for each input
 - Hence it requires the same user index for these 3 inputs simultaneously to ensure correctness
- Accumulator does not consider order of the value stored in it
 - Classical ring signature has user' order

Our solution:

- Add user index to the public keys before accumulating it

Advantage: constant proof size due to the use of accumulator

Disadvantage: trusted setup in accumulator

News about our ESORICS 2017 publication:

👤 kryptomoney ⌚ August 22, 2017 📁 Crypto
Currency, Latest Posts 💬 0

Monero: Cryptocurrency for True Anonymity

Monero was recently in top gainers in the cryptocurrency world, almost doubling from \$50 to \$100 within the last couple of hours and then settling at around \$80. The current surge can be attributed to two major developments in the last month. First, Monero is rumoured to be listed on Bithumb, South Korea's largest bitcoin exchange starting August 23. Second, contributions from the Monero development team in the technical whitepaper, titled "RingCT 2.0: A Compact Linkable Ring Signature Based Protocol for Blockchain Cryptocurrency Monero" by Sun et al which is set to be presented at ESORICS 2017.



BulletRingCT

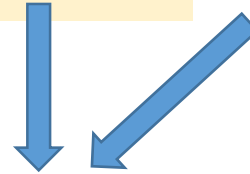
- Current work: RingCT from the Bulletproof framework

BulletRingCT idea:

- Vector commitment $C = h^r y_1^{b_1} y_2^{b_2} \dots y_n^{b_n}$
- If (y_1, y_2, \dots, y_n) are public keys and $\vec{b}_L = (b_1, b_2, \dots, b_n)$ is a **binary vector of Hamming weight 1**, then the commitment C is a commitment to one public key

Bulletproof framework:

- ZK proof for inner product relation
$$\langle \vec{a}, \vec{b} \rangle = x$$
- Proof size is $\log |\vec{a}|$
- No trusted setup



- Define $\vec{b}_R = \vec{1} - \vec{b}_L$
- $\vec{b}_L = (b_1, b_2, \dots, b_n)$ is a **binary vector of Hamming weight 1** can be shown as:
$$\vec{b}_L \circ \vec{b}_R = \vec{0}, \quad \vec{b}_R = \vec{1} - \vec{b}_L, \quad \langle \vec{b}_L, \vec{1} \rangle = 1$$
- We can use Bulletproof for showing inner product with log-size proof!
 - Caution: Bulletproof requires that the DL of y_1, y_2, \dots, y_n are unknown in the security proof.
 - No problem for ring signature
 - Cannot go through the security model of linkable ring signatures! So we have to hash function to formulate the public key set as well.

BulletRingCT

Signature Size

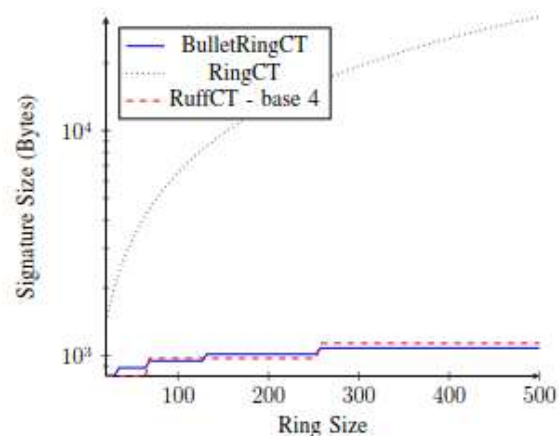


Fig. 2. Comparison of BulletRingCT, RuffCT and RingCT for a transaction with 2 inputs, with small ring size.

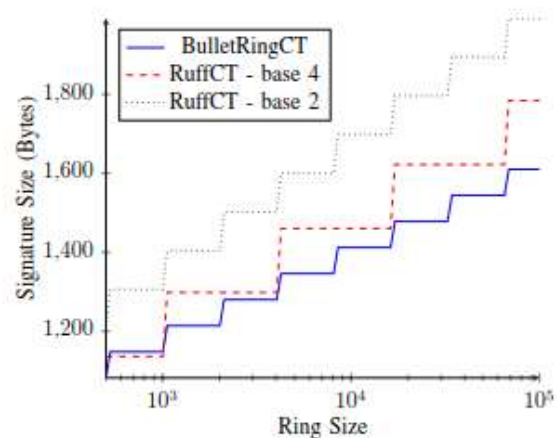


Fig. 3. Comparison of BulletRingCT and RuffCT for a transaction with 2 inputs, with large ring size.

Running Time

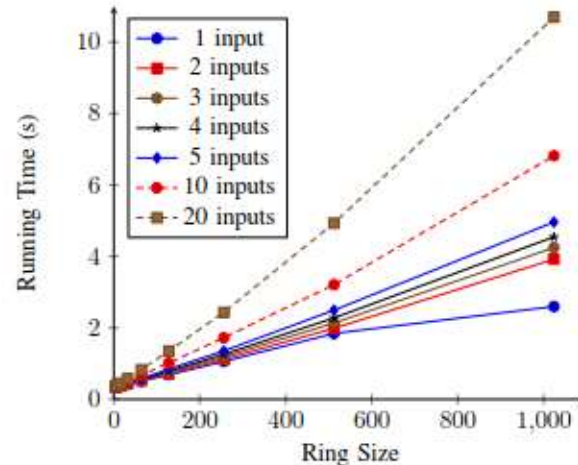


Fig. 4. Running Time of Spend in BulletRingCT.

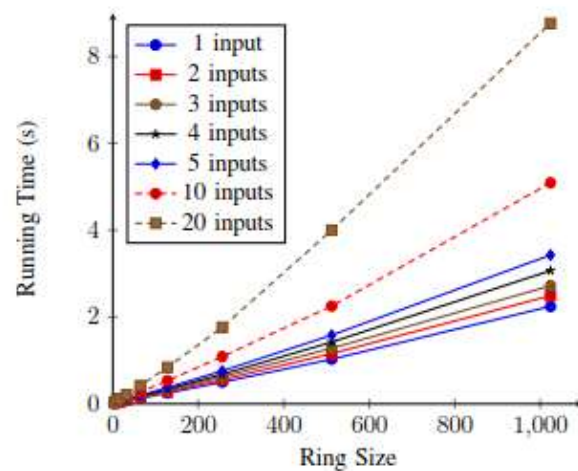


Fig. 5. Running Time of Verify in BulletRingCT.

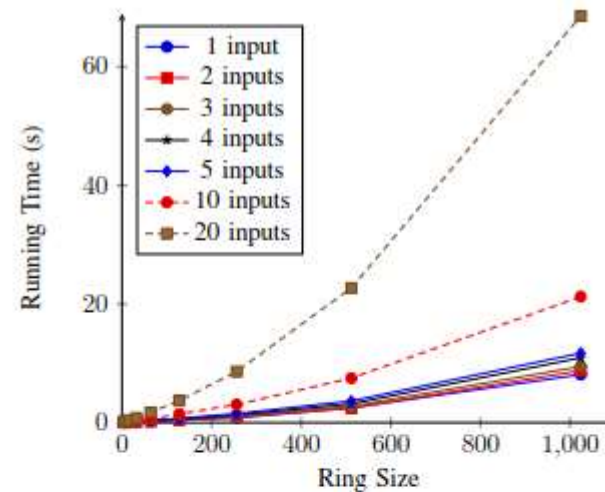


Fig. 6. Running Time of Spend in RuffCT.

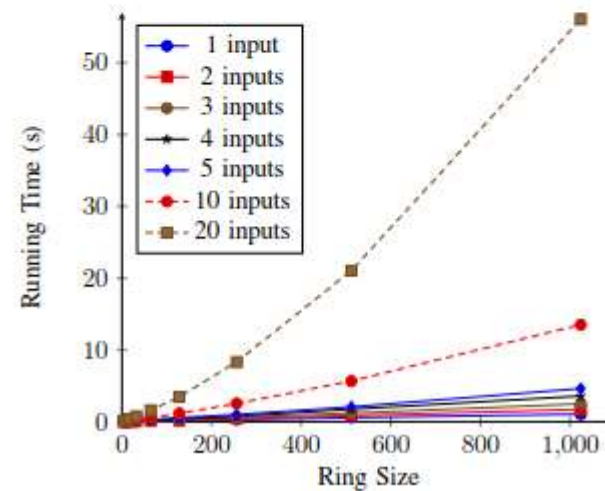
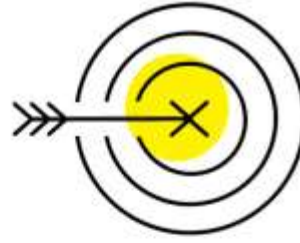


Fig. 7. Running Time of Verify in RuffCT.

Conclusion

- Blockchain is a promising technology, especially for the finance industry.
 - Other killer applications to be determined
- Many active research areas:
 - Short term: Consensus, privacy
 - Medium-Long term: DAG data structure, inter-chain transactions, smart contract security, sharding, layer 2, cryptoeconomics...
- Ring signature is one important tools for privacy in blockchain applications.

Top 10 Strategic Technology Trends for 2018



Intelligent



AI Foundations



Intelligent Apps
and Analytics



Intelligent Things



Digital



Digital Twins



Cloud to the Edge



Conversational
Platform



Immersive
Experience



Mesh



Blockchain



Event-Driven



Continuous Adaptive
Risk and Trust

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